

Classifying the computational power of stochastic physical oracles

Edwin Beggs¹, Pedro Cortez², José Félix Costa^{2,3}, and John V. Tucker¹

¹ College of Science, Swansea University, Singleton Park, Swansea, SA2 8PP, Wales, U.K.

² Department of Mathematics, Instituto Superior Técnico, Universidade de Lisboa, Lisboa, Portugal

³ Centro de Filosofia das Ciências da Universidade de Lisboa, Lisboa, Portugal

Consider an algorithm requesting information from an external source – an *oracle* – the terminology originates with Alan Turing [7]. Emil Post [6] used oracles to study computability.

However, suppose the external source is not a pure mathematical entity but a *physical device or environment*. Suppose the requests are for measurements of physical quantities. We call this external source a *physical oracle*. Algorithms with physical oracles may be found in measurement experiments, and in controlling machines. We ask: *What is the computational power of adding a physical oracle? How does the computational theory depend upon the physical theories and models?*

In [2,3] we developed a computability and complexity theory for physical oracles. The computational classification needed non-uniform complexity classes [1], especially P/\log^* and $BPP//\log^*$ [5]. Using case studies, we formulated axioms expressing properties common to wide classes of physical systems [4].

Here we review physical oracles and report new results broadening their scope by using *non-deterministic physical systems*. Physical oracles with probabilistic theories we call *stochastic physical oracles*. We examine examples of three types of non-deterministic systems, those that are physically nondeterministic, as in quantum phenomena; physically deterministic but whose physical theory is non-deterministic, as in statistical mechanics; and physically deterministic but whose computational theory is non-deterministic caused by error margins. We prove:

Theorem 1. *Let SPO be the axioms for stochastic physical oracles. Let P be a physical system whose behaviour depends upon a physical quantity or parameter σ . Suppose P satisfies the axioms of SPO. Then: a set $A \subset \{0, 1\}^*$ is decidable in polynomial time by a Turing machine with physical oracle P and unknown parameter σ if, and only if, $A \in BPP//\log^*$.*

References

1. José Luis Balcázar, Josep Díaz, and Joaquim Gabarró. *Structural Complexity I*. Springer-Verlag, 2nd edition, 1988, 1995.

2. Edwin Beggs, José Félix Costa, Bruno Loff, & John V. Tucker. Computational complexity with experiments as oracles. *Proc. Royal Soc., A (Math., Phys. and Eng. Sci.)*, 464(2098):2777–2801, 2008.
3. Edwin Beggs, José Félix Costa, Bruno Loff, and John V. Tucker. Computational complexity with experiments as oracles II. Upper bounds. *Proc. Royal Soc., Ser. A*, 465(2105):1453–1465, 2009.
4. Edwin J. Beggs, José Félix Costa, and John V. Tucker. Axiomatizing physical experiments as oracles to algorithms. *Phil. Trans. Royal Soc., Ser. A (Math., Phys. and Eng. Sci.)*, 370(12):3359–3384, 2012.
5. Edwin J. Beggs, José Félix Costa, Diogo Poças, and John V. Tucker. An analogue-digital Church-Turing thesis. *Int. Journal of Foundations of Computer Science*, 25(4):373–390, 2014.
6. Emil Post. Degrees of recursive unsolvability. *Bull. of the American Math. Society*, 54, 1948, 641–642.
7. Alan Turing. Systems of logic based on ordinals. *Proceedings of the London Mathematical Society*, Second series, 45:161–228, 1939